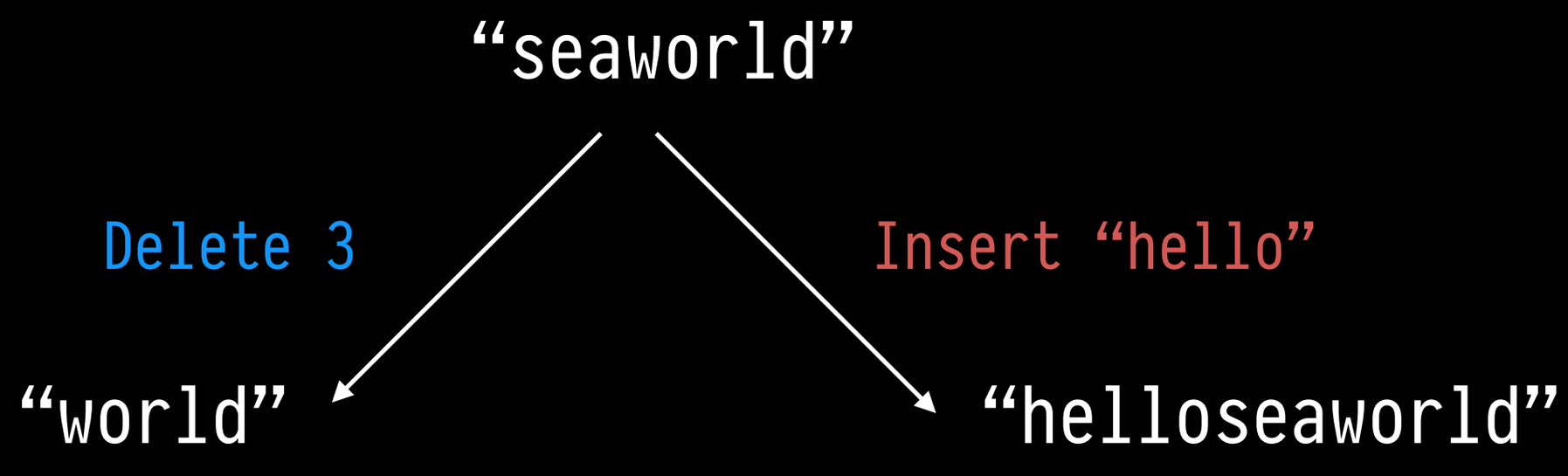
Streaming OT

or “the wonders of Gabriel Gonzalez”

OT

```
data Op = Retain Int | Delete Int | Insert String
type Delta = [Op]
      ^a diff
```

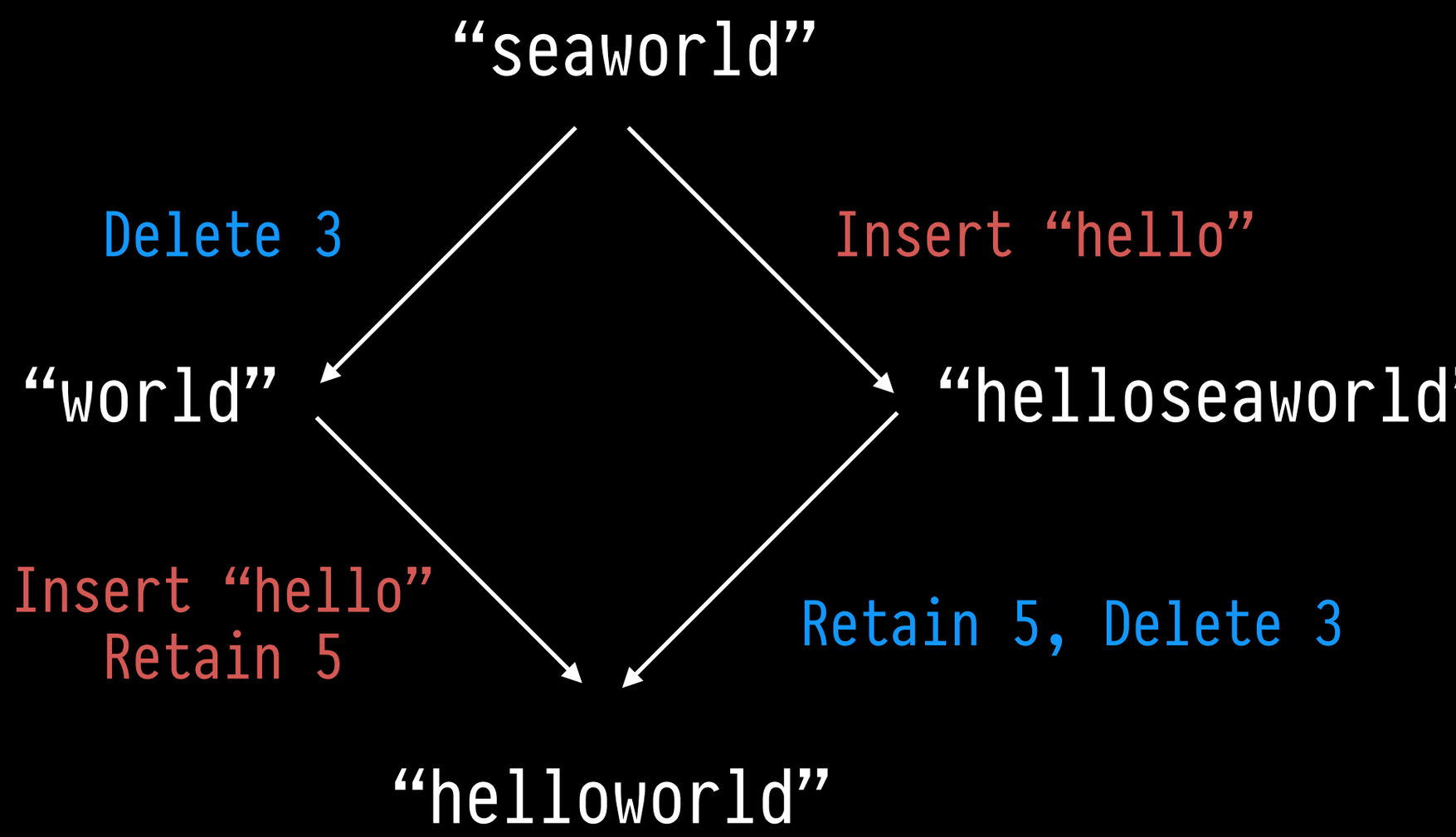
OT



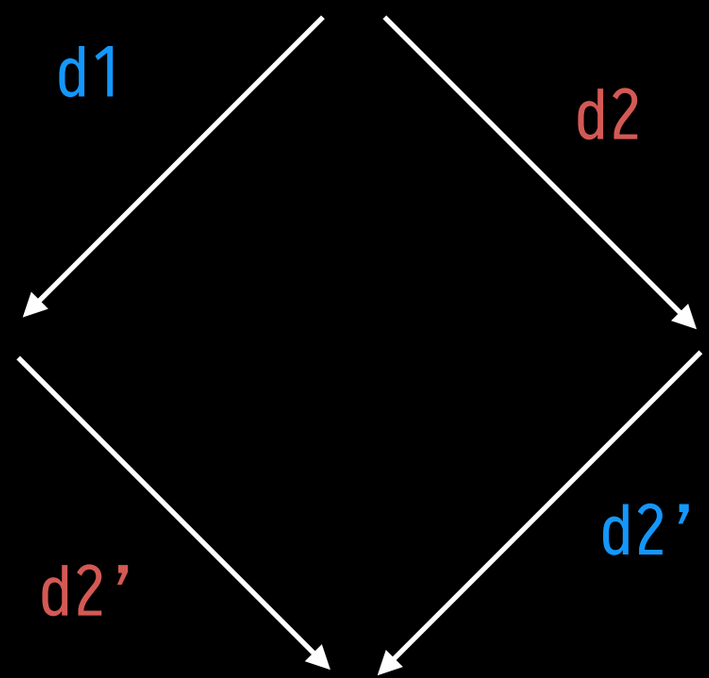
OT

```
                incoming op    history
                v
transform (Delete 3:op1s) (Insert “hello”:op2s)
= [ Insert “hello”
  , Retain 5
  , transform (Delete 3:op1s) op2s
```

OT


“seaworld”
Delete 3
Insert “hello”
“world”
“helloseaworld”
Insert “hello”
Retain 5
Retain 5, Delete 3
“helloworld”

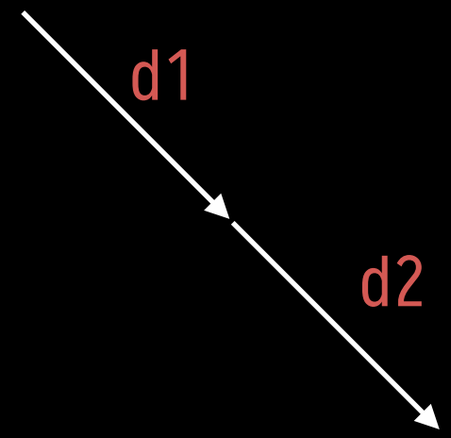
OT




```
d1’ = transform d1 against d2
d2’ = transform d2 against d1

d1 <> d1’ == d2 <> d2’
```

proof: easy induction and case distinction. QED.

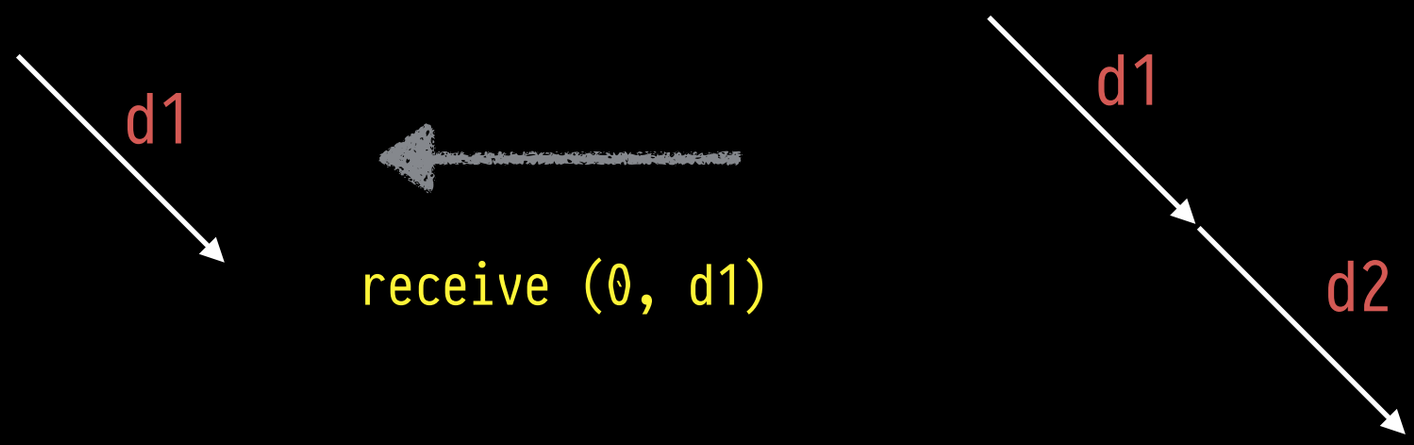
OT



Client

Server

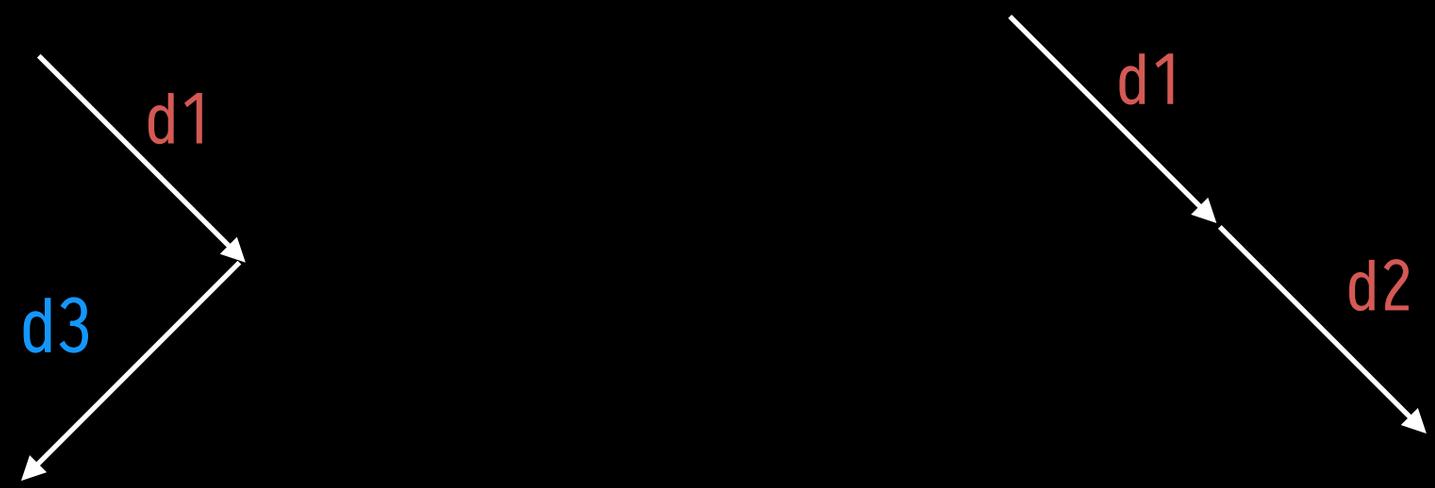
OT



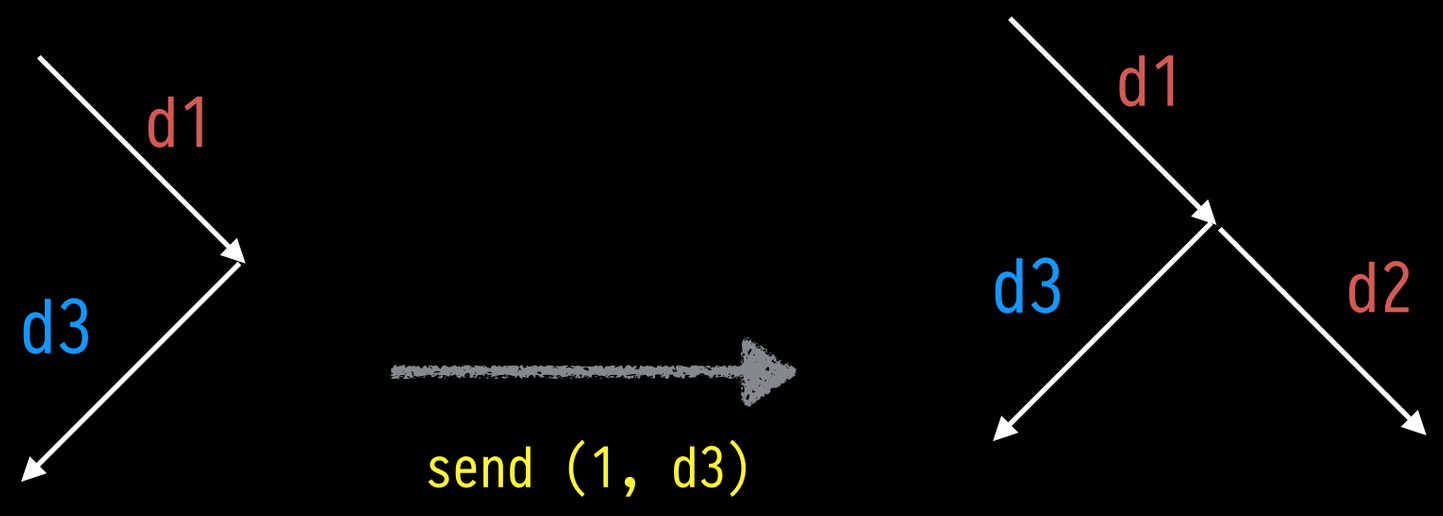
Client

Server

OT



OT

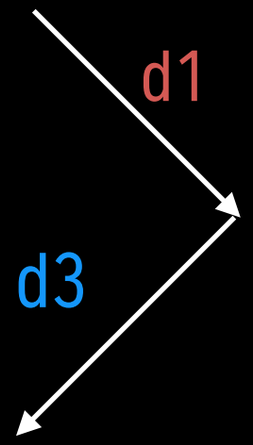


Client

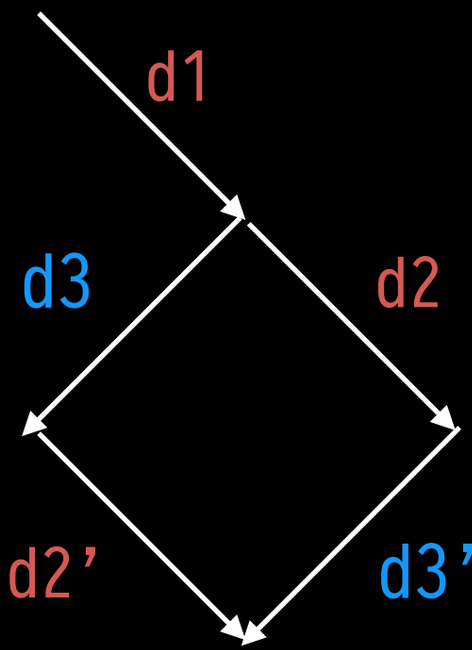
Server

OT

OT



Client

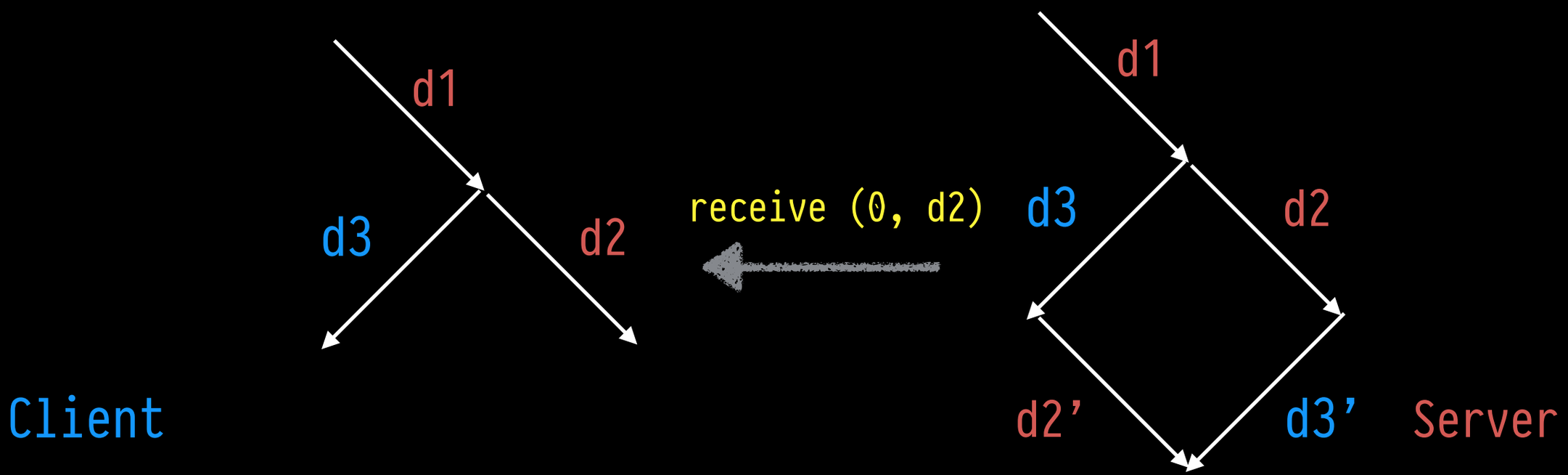


Server

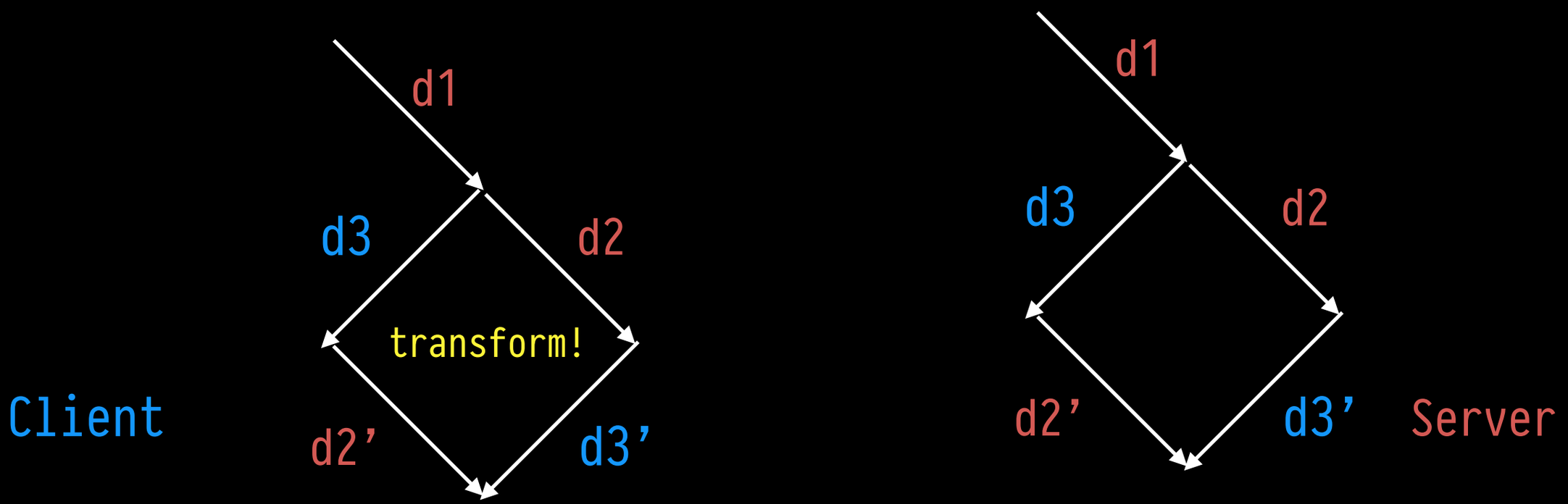
sends d3’ to other clients

new history: [d2’]

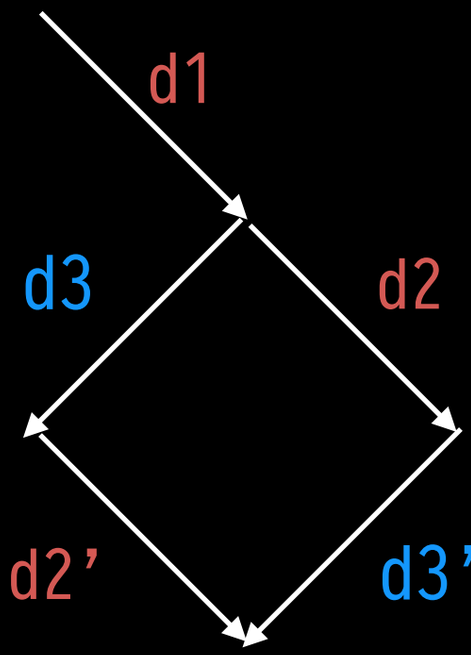
OT



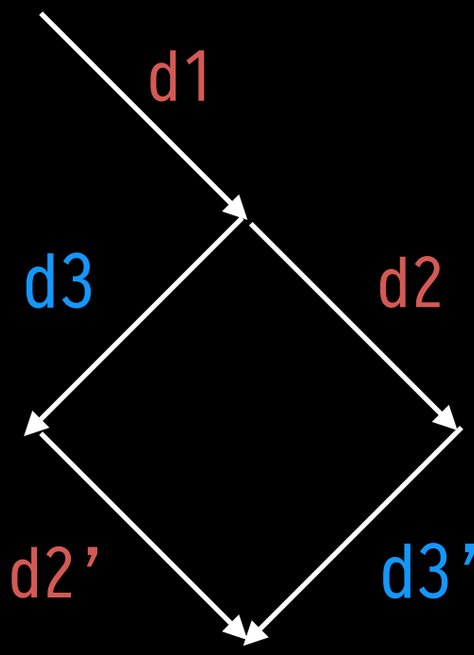
OT



OT



Client

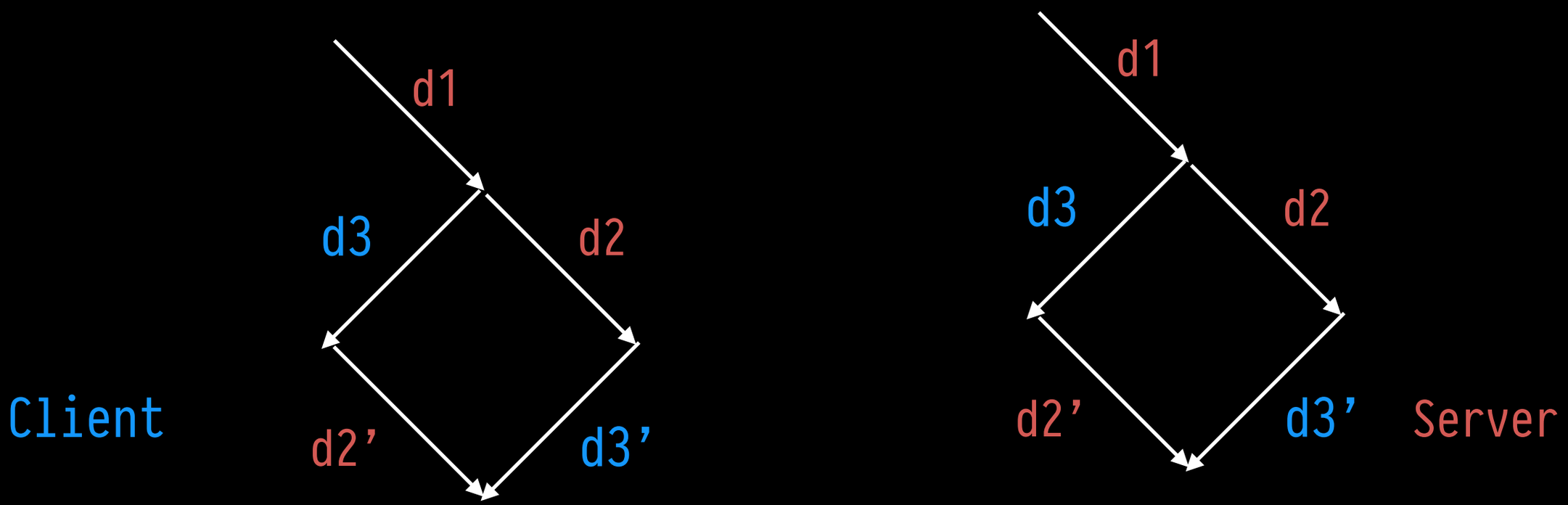


Server

apply d2’ to document

new history: [d3’]

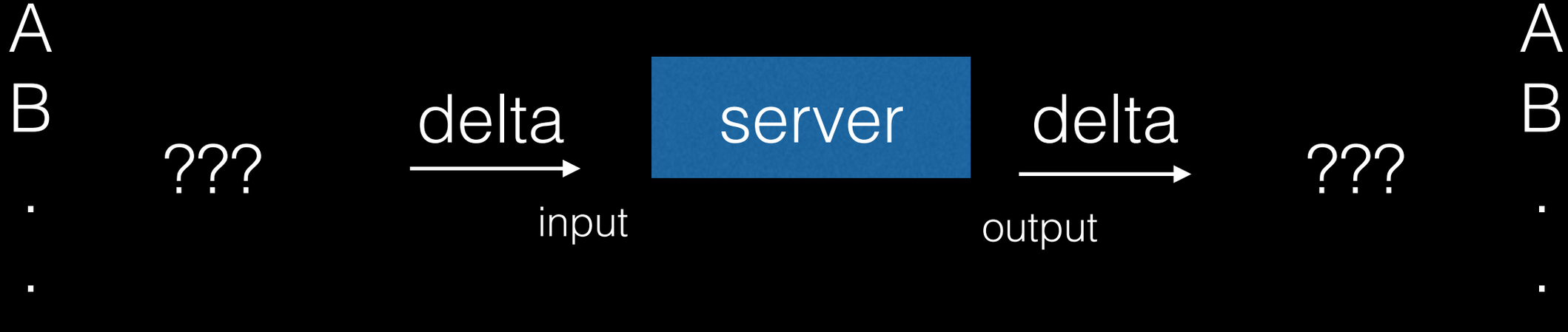
OT



d1 <> d3 <> d2’ == d1 <> d2 <> d3’

by transform property

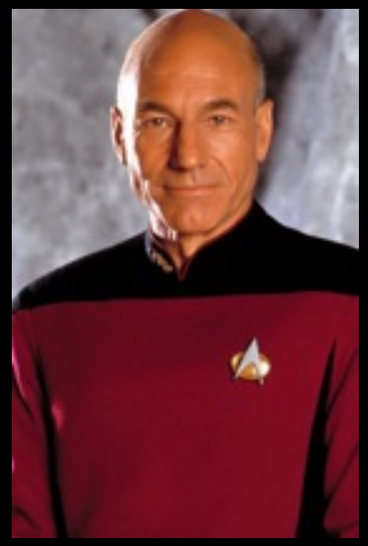
Pipes



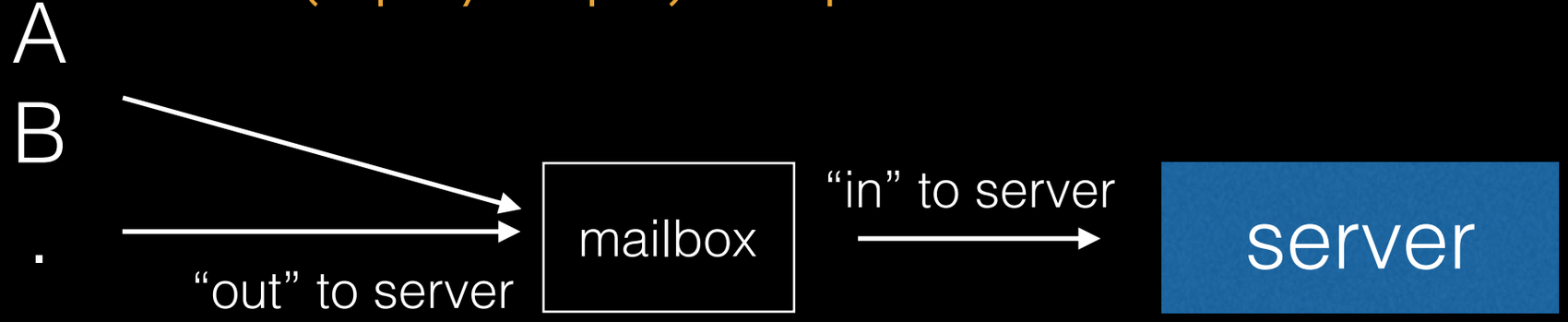

```
server :: Pipe (Signed TrackedDelta)
               (Signed Delta)
               (StateT (Map ClientID [Delta]) m)
```

it’s pure, woo!

Actors!
(more like pi-calculus)



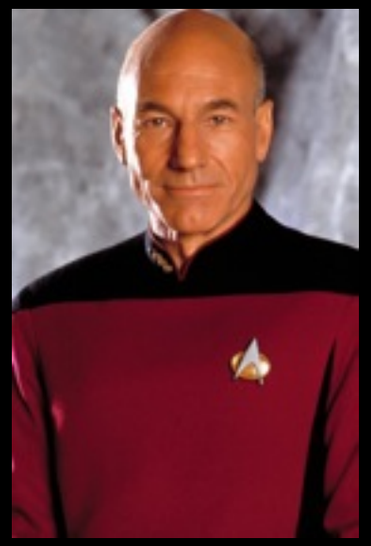
```
spawn :: Buffer a -> IO (Output a, Input a)

(input, output) <- spawn Unbounded
```




```
fromInput input :: Producer (Signed TrackedDelta)

fromInput input >-> server >-> . . .
```

streaming input, woo!

Actors!



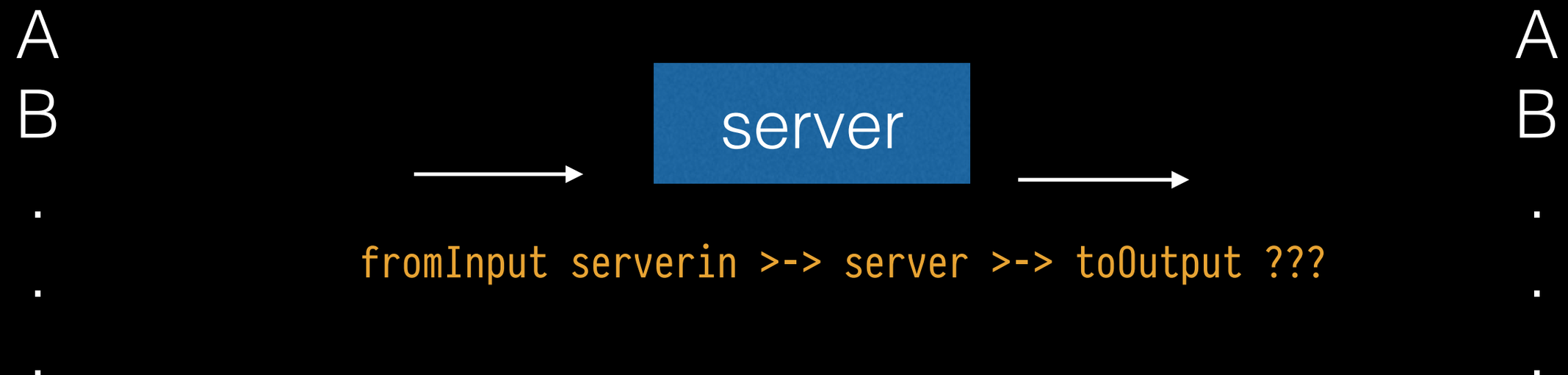
```
spawn :: Buffer a -> IO (Output a, Input a)

(input, output) <- spawn Unbounded
```


```
reader id tcp :: Producer (Signed TrackedDelta) m ()
reader id tcp =   P.fromHandle tcp
              >-> P.read
              >-> P.map (Signed id)
```

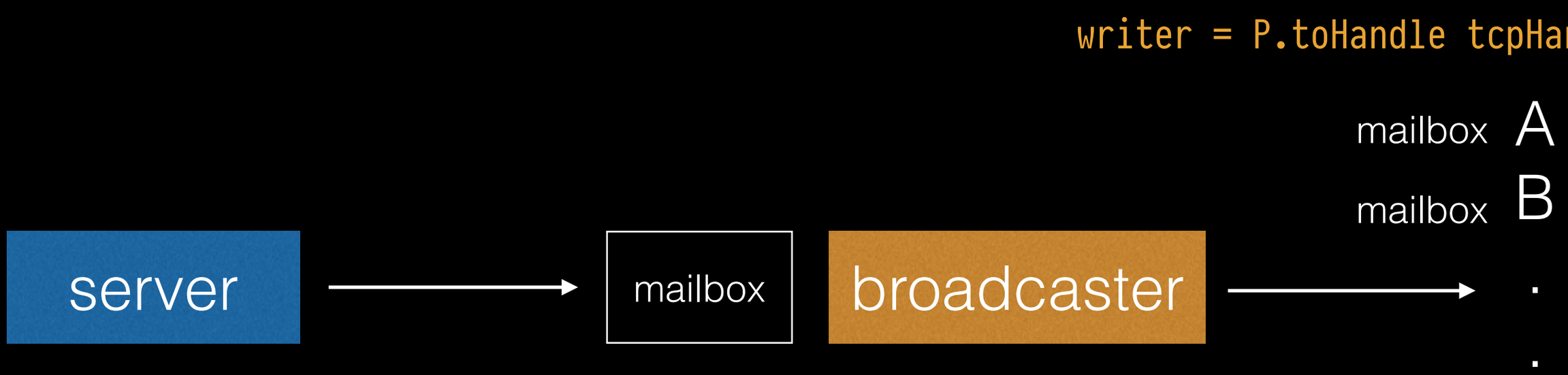
```
reader clientID handle >-> toOutput output
```

Pipes

```
    fromHandle
>-> …
>-> toOutput serverout
```

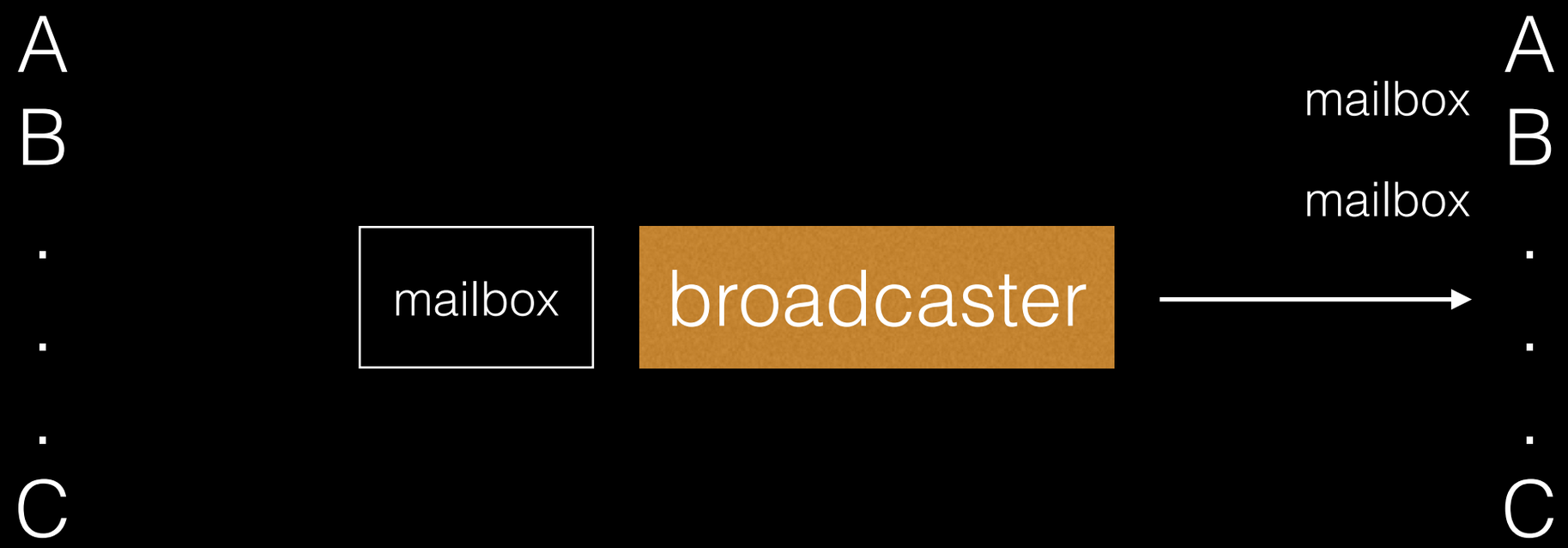


Pipes




```
broadcast :: Consumer BroadcastMsg
                      (StateT (Map ClientID
                                   ( Int
                                     ^ no. deltas client reported to have seen
                                   , [Output TrackedDelta]
                                     ^ client writer's mailbox)
                               m)
```

Pipes



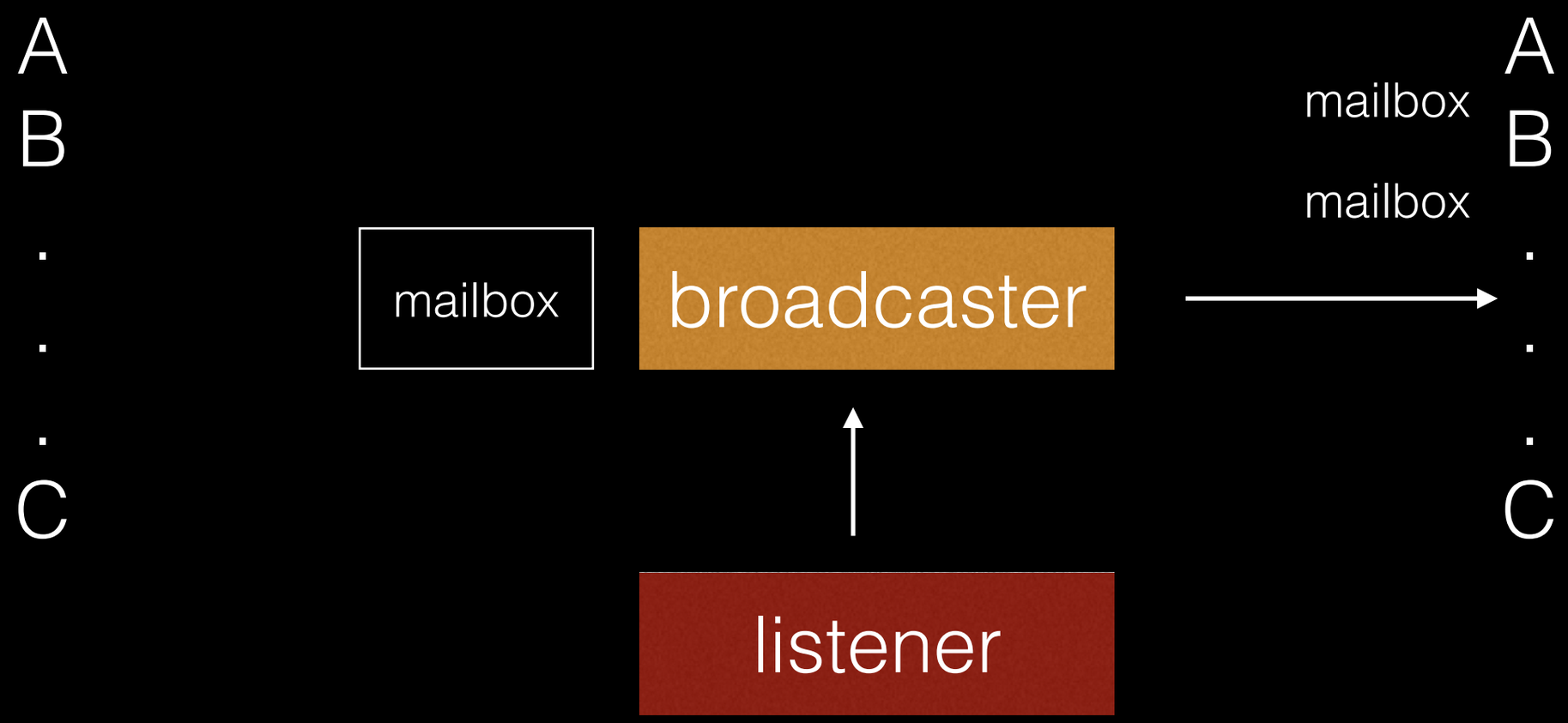
listener

listen for new clients
start a reader and writer

Pipes


```
(writerout, writerin, writerseal) <- spawn' Unbounded
writerid <- forkIO $ catch (runEffect $ fromInput writerin >-> writer h)
                           (\Disconnect -> atomically $ writerseal)
forkIO $ do runEffect $ reader cid h >-> toOutput readerout
            throwTo writerid Disconnect
return $ Client cid writerout
```

Pipes




```
listener serverout handle >-> toOutput broadcastout
```

Client

model

```
Model ( Int       – no of ops to tell server to drop from history
      , [Delta])  - ops sent, not confirmed by server
      From To


  data From = FServer TrackedDelta    – op from server
            | FUser   Delta           - op from editor

  data To   = TServer TrackedDelta
            | TUser   Delta
```

MVC (no not that one)

editor controller

editor view

```
model      :: Model (Int, [Delta]) From To
controller :: Managed (Controller From)
view       :: Managed (View To)
```

MVC (no not that one)

editor controller

editor view

```
controller tcp ws
=  fmap (fmap FServer) (controllerServer tcp)
<> fmap (fmap FUser)   (controllerUser   ws)

controllerServer = M.producer Single (P.fromHandle tcp >-> P.read)
^ deltas from server

controllerUser   = M.producer Single (void $ L.view P.decoded $ fromWS ws)
^ deltas from JSON from the js-based editor (yuck)
```

MVC bug

```
controller = M.producer Single $ P.stdinLn
view       = M.consumer $ P.stdoutLn
```

```
a
b
<a>
c
<b>
```

So

- pipes are strongly principled abstractions
- good ecosystem
- mvc getting there
- Gabriel is great